

An Implementation of a Software UART Using the TMS320C25

APPLICATION REPORT: SPRA077

*David Zalac
Digital Signal Processor Products-
Semiconductor Group
Texas Instruments*

Digital Signal Processing Solutions



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CONTACT INFORMATION

US TMS320 HOTLINE	(281) 274-2320
US TMS320 FAX	(281) 274-2324
US TMS320 BBS	(281) 274-2323
US TMS320 email	dsph@ti.com

An Implementation of a Software UART Using the TMS320C25

Abstract

When transmitting to and receiving data from a processing engine such as the TMS320C25 digital signal processor, it is common to encounter problems when interfacing to asynchronous devices. This chapter provides a look at software implementation of a Universal Asynchronous Receiver and Transmitter (UART) which enables communication with asynchronous serial devices with a minimum of external hardware. Topics covered include:

- ❑ Characteristics of asynchronous communications
- ❑ Limitations of and reasons where use of a UART are justifiable
- ❑ An overview of the functions provided by a UART
- ❑ UART Implementation
- ❑ A list of References
- ❑ An appendix of source code

Supporting figures illustrate the following:

- ❑ UART word format and architecture
- ❑ Timer interrupt service routine
- ❑ Transmitter routine
- ❑ UART status word
- ❑ Receiver routine
- ❑ RS-232 interface



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Introduction

Interfacing to asynchronous devices is a common problem in transmitting to and receiving data from a processing engine such as the TMS320C25 digital signal processor. This report describes a software implementation of a Universal Asynchronous Receiver & Transmitter (UART) that provides the ability to communicate with asynchronous serial devices in a system with a minimum of external hardware.

Asynchronous communications are characterized by the absence of a timing reference such as a clock or framing signal. Various tradeoffs arise from this distinction from synchronous communications in terms of hardware and software requirements and data throughput capacity. Synchronous communications require a timing reference, but otherwise have minimal hardware and software requirements. Asynchronous communications require a mechanism for deriving a timing reference from the received signal. Additionally, various error-checking functions are typically implemented. These requirements impose hardware and/or software overhead that is not imposed in the synchronous case. Moreover, synchronous interfaces can typically support much higher data throughput rates than asynchronous interfaces.

Implementing a UART in software imposes CPU overhead whose acceptability is application-dependent. In applications where the overall data throughput rate is sufficiently low, or in cases in which a UART is to be used only for booting system memory at powerup, use of a software UART may be justifiable. A hardware solution (i.e., a UART IC) may be more appropriate in high data rate applications and in applications requiring low I/O overhead. A detailed analysis of overhead imposed by the TMS320C25 software UART is given later in this report.

A high-speed synchronous serial interface is provided by the on-chip serial port of the TMS320C25. A full description and specification of the serial port may be found in the *Second-Generation TMS320 User's Guide*. [4]

Overview

The functions provided by a UART are simply the transmission and reception of serial data and the checking and signalling of various error conditions. These functions are described in detail in the following sections.

Data Format

Shown in Figure 1 is the layout of a word in a format assumed by the UART. Bit 0 is a space (logic low) and is referred to as the start bit. Bits 1 through N are the N data bits of the word with the LSB occupying bit position 1. Typically, N has a value of 5, 6, 7, or 8. The maximum value of N is given by $N_{max} = 14 - M$, where M is the

number of stop bits. Bit N+1 is referred to as the parity bit and has a value such that the total number of ones in the word (bits 1 through N+1) is odd if odd parity is selected and even if even parity is selected. Bits N+2 through N+M+1 are referred to as stop bits, and each has a value of one. The total word length WORD_LEN is thus given by $WORD_LEN = N+M+2$.

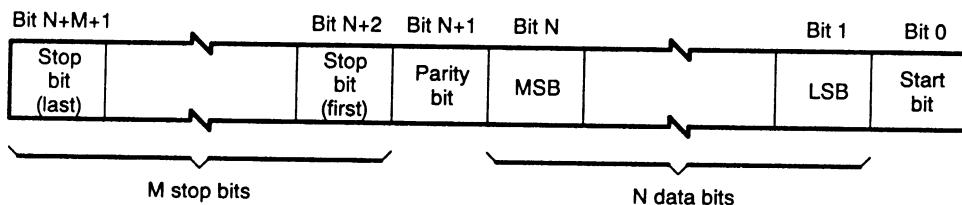


Figure 1. UART Word Format

Data Reception

Reception of a data word starts with detection of the start bit. One way of performing start bit detection is to sample the input data signal at a rate that is large compared to the bit rate, then testing each sample for a space (logic low). An optional check can be performed to verify that the first logic low detected represents a valid start bit and not just noise. This check is performed by testing that the input signal is low one-half-bit duration after the start bit has been detected.

Once the start bit has been detected, the UART simply recovers the data from the input signal and keeps track of the data parity. The parity is checked against the received parity bit after all the data bits have been received. Finally, the integrity of the word framing is checked by testing that the input signal is high when the first stop bit is expected.

Data Transmission

Transmission is considerably simpler than reception in that timing information does not have to be recovered from an asynchronous signal. Furthermore, no error checking is performed by the UART transmitter. Transmitting a data word is preceded by appropriately formatting the data to be transmitted; i.e., adding start, stop and parity bits. Formatting is done in TMS320C25 software. The output signal is generated from the data and appears on the UART's output signal line.

Implementation

The UART implementation described in this report makes use of two TMS320C25 general-purpose I/O pins (XF and BIO/) and the timer interrupt. The input signal is received on the BIO pin via the TMS320C25 BIOZ instruction. The output signal appears on the general-purpose flag pin XF. The state of XF is controlled in software via the SXF and RXF instructions. The TMS320C25 serial port is not used. As shown in Figure 2,

the transmitter and receiver are “serviced” each time a timer interrupt is generated. The timer interrupt rate is an integer multiple (K) of the bit rate.

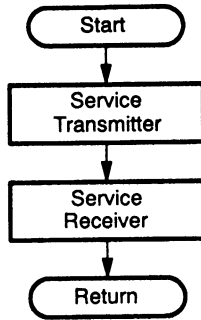


Figure 2. Timer Interrupt Service Routine

Several pieces of code comprise the UART software:

1. Timer interrupt service routine
2. UART__INIT initialization routine
3. XMT routine
4. RCV routine
5. PUT__DATA
6. GET__DATA
7. XCOMPOSE

The UART transmitter and receiver are located in the timer interrupt service routine. (No context save/restore is included in this interrupt service routine. Refer to the “Precautions” section for details.) UART__INIT initializes the UART with the values appearing in the assembly-time constants section of the source listing. XMT and RCV are user-written routines that interface the UART to the user’s program. XCOMPOSE, PUT__DATA and GET__DATA are auxiliary routines available to the user for executing UART interface housekeeping tasks. Each of these seven routines is described in detail later in this report.

Figure 3 shows the UART structure and the transmit and receive data paths. The transmit and receive buffers TDATA and RDATA are 16-bit wide TMS320C25 on-chip data memory locations. The transmit and receive software shift registers TSHF and RSHF are also located in data memory. The UART status word USTAT maintains UART status and error information. USTAT will be discussed in detail later in this report. The value of USTAT is written to I/O port UARTPORT each time USTAT is updated, thus allowing the capability of externally monitoring the UART status. The input and output pins ($\overline{\text{BIO}}$ and XF) may be interfaced to RS-232-compatible transmit and receive lines. Finally, the locations in data memory for transmit and receive data are pointed at by two TMS320C25 auxiliary registers AR(OPT_PTR) and AR(INP_PTR), respectively.

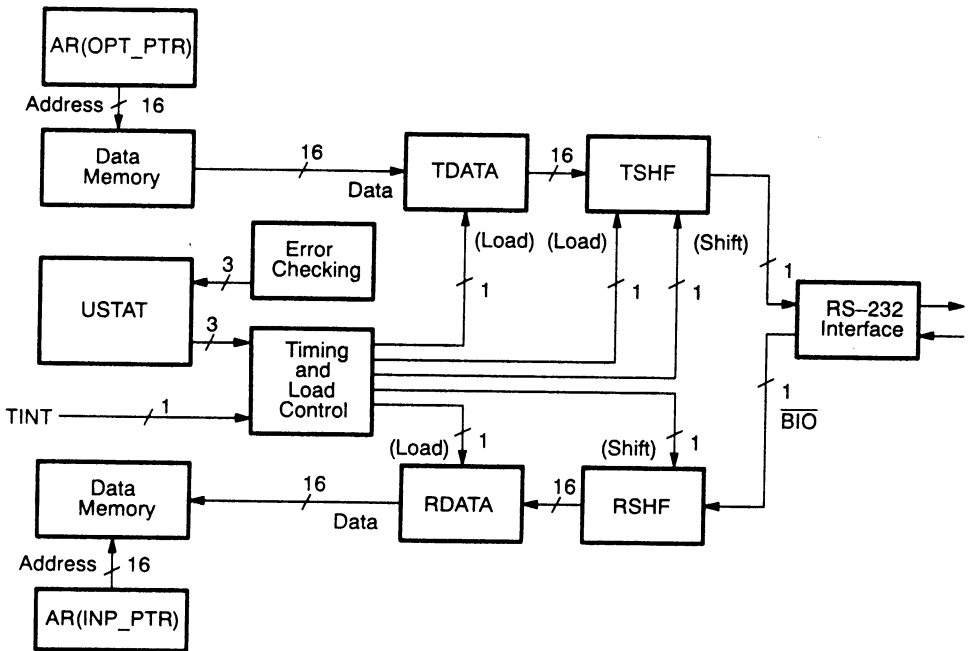


Figure 3. UART Architecture

All UART variables are mapped into TMS320C25 on-chip RAM block B2. The code size is 332 words and can be executed from on-chip ROM, EPROM (TMS320E25), or off-chip program memory. The maximum bit rate supported is 19.2 kilobits per second (full duplex).

Transmitter

Figure 4 shows a flowchart of the transmitter routine. When the transmitter completes transmission of the current word, a new word (TDATA) is loaded into the transmitter software shift register if the TDA (Transmit Data Available) flag is set.

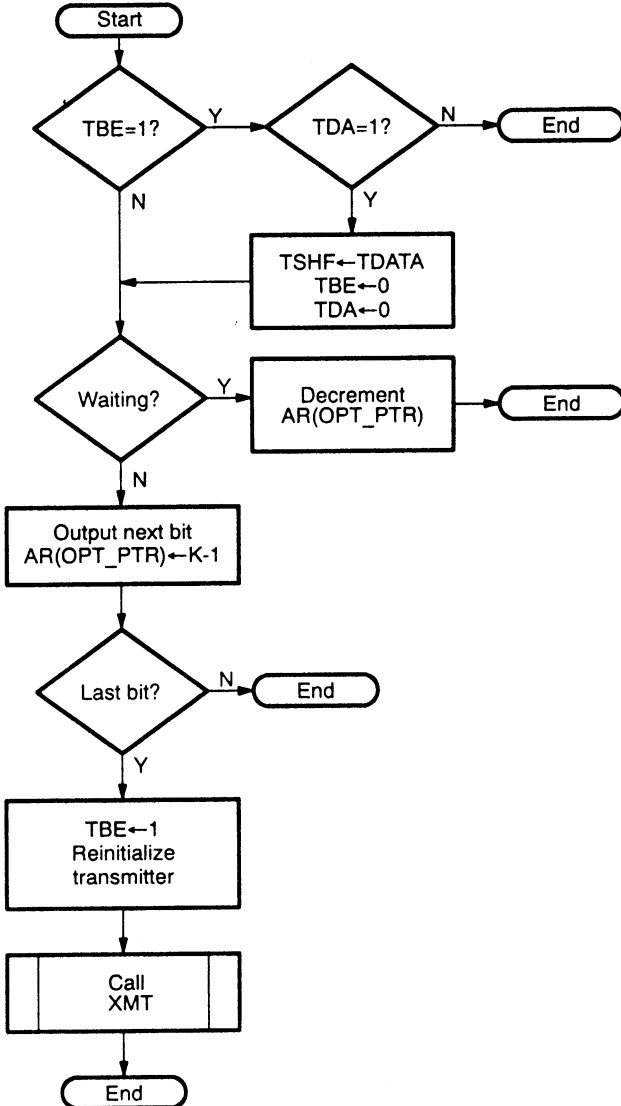


Figure 4. Transmitter Routine

TDA is one of 6 flags residing in the UART status word (see Figure 5). The word to be transmitted is shifted out on the XF pin at the user-specified bit rate.

Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
Not used	Not used	TDA	RDE	RDA	ROR	FRM	RPE

Figure 5. UART Status Word

Feeding data to the transmitter consists of three steps: appropriately reformatting the data to be transmitted (i.e., adding start, parity, and stop bits), loading the data into the UART variable TDATA, and indicating to the UART that valid transmit data is present in TDATA by setting the TDA flag. The first step can be accomplished by pointing AR(OPT_PTR) at the data to be transmitted and calling auxiliary routine XCOMPOSE. XCOMPOSE does an in-place reformatting of the data per the values in the UART parameter variables (see source listing). Routine PUT_DATA may be called to load TDATA and set TDA.

Each time the transmit shift register empties, a call to XMT is made. Management of AR(OPT_PTR) and calls to XCOMPOSE and PUT_DATA may be made from XMT. Alternatively, these functions can be executed from the user's program. This is the preferred approach, because all code in XMT adds to the maximum path length through the timer interrupt routine and thus decreases the maximum bit rate. However, calls to PUT_DATA from the user's program should be made only if TDA = 0. If this condition is not satisfied, the current word to be transmitted will be overwritten. (The condition TDA = 0 is guaranteed if PUT_DATA calls are made from XMT and needn't be checked.)

Initiation of transmission of the first word in a string of words (string = one or more words) must be made from the user's program by calling XCOMPOSE and PUT_DATA, as no XMT calls can be made until the transmitter is started.

The status of the transmitter can be ascertained by reading bit 5 of the UART status word, as shown in Figure 5:

USTAT

BIT 5 Transmit Data Available TDA = 1 indicates to the UART that valid transmit data is present in TDATA.

Receiver

Shown in Figure 6 is a flowchart of the receiver routine. The state of the receiver is indicated by the value of RSTAT and bits 0 - 4 of the UART status word as shown below.

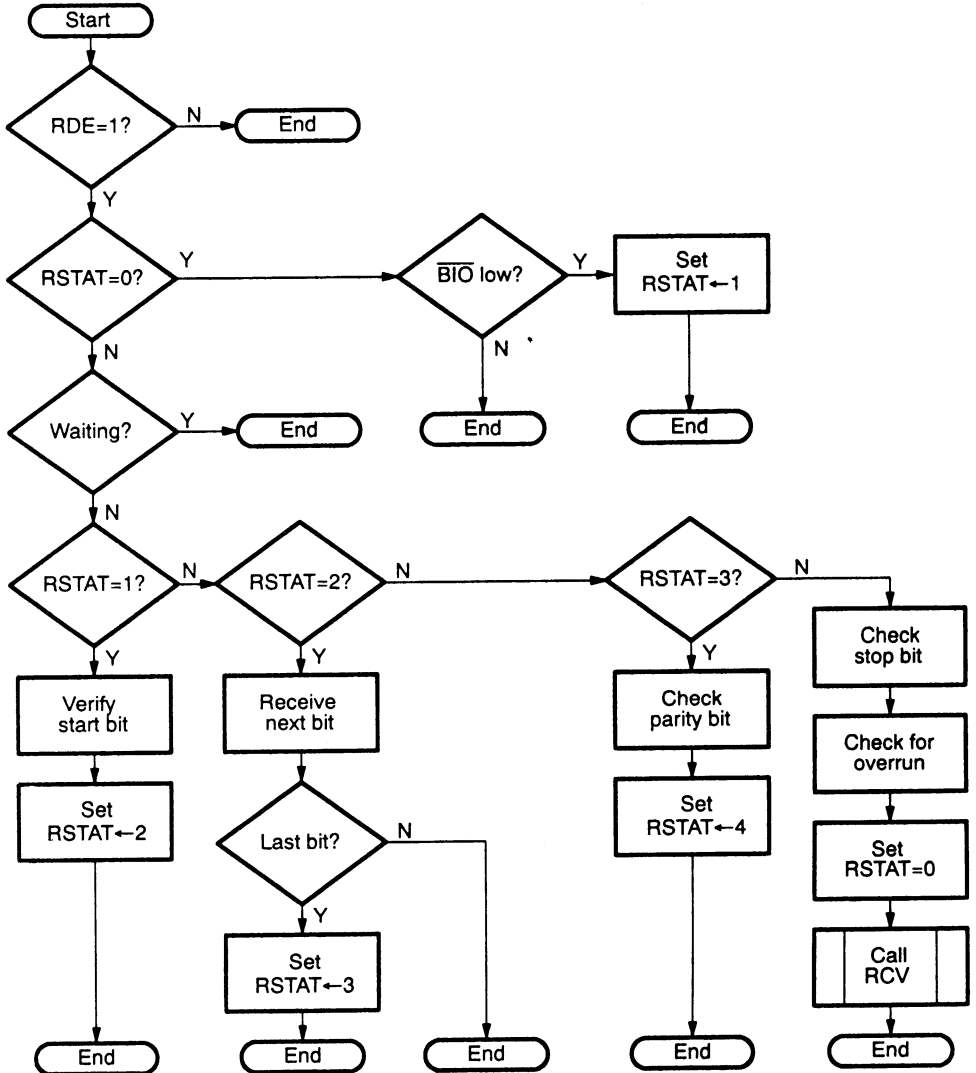


Figure 6. Receiver Routine

RSTAT Receiver Status

0	Waiting for start bit
1	Waiting for start bit center
2	Waiting for data bit
3	Waiting for parity bit
4	Waiting for stop bit

USTAT

BIT 0	Receive Parity Error	If receive parity checking is active (RPACTIVE = 1), RPE is set if a parity error is detected.
BIT 1	Framing Error	This bit is set if a logic low is sensed on $\overline{\text{BIO}}$ when the first stop bit is expected.
BIT 2	Receiver Overrun	This bit is set if RDA is not cleared before reception of the next word is completed.
BIT 3	Receive Data Available	This bit is set when reception of a word is completed and indicates the presence of valid data in RDATA.
BIT 4	Receive Data Enable	The receiver routine is bypassed if this bit is a zero.

Reception is initiated by setting RDE. Each time a complete word is received, a call to RCV is made. RCV can call GET_DATA to copy the new data to the location pointed at by AR(INP_PTR) and to clear RDA. If RDA is indeed cleared by every call to RCV, some overhead can be eliminated by deleting the setting, clearing, and checking of RDA because the overrun detect function is superfluous in this case.

The received data in RDATA is right-justified with the LSB in the zero'th bit position and with start, stop, and parity bits stripped.

Overhead and Optimization

The overhead imposed by the UART is primarily determined by the length of the timer interrupt service routine and the timer interrupt rate ($K \times$ bit rate). An expression for overhead is given below:

$$\text{overhead (\%)} = K \times (\text{bit rate}) \times (T + R) \times T_{c(C)} \times 100\%$$

$K = (\text{timer interrupt rate})/(\text{bit rate})$. $T_{c(C)}$ is the period of CLKOUT1 and CLKOUT2. In the timer interrupt service routine shown in the source listing, T and R are given by

$$T = (8 \times \text{WORD_LEN} \times K + 14 \times \text{WORD_LEN} + 15)/(\text{WORD_LEN} \times K)$$

$$R = (17 \times \text{WORD_LEN} \times K + 36 \times \text{WORD_LEN} + 169)/(\text{WORD_LEN} \times K)$$

The values of T and R at WORD_LEN = 10 and K = 16 are T = 9 cycles per timer interrupt and R = 20 cycles per timer interrupt. (These are the correct values when the timer interrupt service routine is in zero-waitstate external program memory. T and R have smaller values if the service routine is in internal program memory.)

T and R represent the average path lengths (in processor cycles) through the transmitter and receiver routines, respectively, in continuous full-duplex operation. Values of overhead for several bit rates and values of K are tabulated in Table 1. Overhead associated with XMT, RCV, GET_DATA, PUT_DATA, and XCOMPOSE code is not included. Note that continuous full-duplex operation constitutes worst-case scenario, a scenario unlikely in applications using asynchronous I/O.

Table 1. UART Overhead (%) vs. Bit Rate (r in bps) and K for WORD_LEN=10†

K/r	300	1200	2400	4800	9600	19,200
6	1	3	5	10	21	42
8	1	3	6	13	26	51
10	1	4	8	15	31	—
12	1	4	9	18	35	—
14	1	5	10	20	40	—
16	1	6	11	22	45	—

†Overhead imposed by the UART is relatively insensitive to WORD_LEN.

There are several ways the user can modify the UART code to reduce the values of T and/or R. Some of these involve eliminating the setting, clearing, and checking of flags in the USTAT register that are not necessary in a fixed configuration. Others involve streamlining of the interface between the UART software and the user's program.

UART Configuration

The initial values of the UART parameters (e.g., bit rate and parity type) appear in the assembly-time constants section of the source listing given in the appendix of this report. Note that an initial-value constant exists for each UART parameter and has the same name as its corresponding parameter, but with an 'I' prefix. If the UART is to be run in a fixed configuration, the user needs only to modify the initial value parameters and re-assemble and link the program. However, if the configuration is to be modified "on the fly", the following measures need to be taken:

To Respecify	Modify
RPACTIVE	RPACTIVE (0 or 1)
K	KM1, K2M1, and TINTPER (KM1 = K-1; K2M1 = K/2-1)
N	N and WORDLEN, CALL PARINIT
WORD_LEN	WORD_LEN
ODD	ODD, CALL PARINIT (initialize parity templates)
TINTPER	TINTPER
# of stop bits	WORD_LEN (# of stop bits = WORD_LEN-N-2)

Precautions

No context save/restore is provided in the timer interrupt service routine as the user will want to write and optimize this part of the routine for his own application. The timer interrupt routine affects the following registers and memory locations:

CPU Registers	Memory Locations
Accumulator	60h-77h (RAM block B2)
T register	dma<AR(INP_PTR)>
P register	TIM register
Auxiliary registers 1-7	PRD register
Status regs ST0 & ST1	

The timer interrupt routine uses two levels of stack plus as many levels as are required to accommodate subroutine calls from XMT and RCV.

If the PRD register contains a value less than 64 (19.2 kbps @ K = 8 or 9.6 kbps @ K = 16), the sampling of some receive bits may be significantly delayed from bit interval centers and some transmit signal edges may be delayed.

The actual transmit and receiver-sampling bit rate r is given by

$$r = 1/[P \times K \times T_{c(C)}],$$

where P is the sum of the contents of the PRD register and one. If no integer value of P exists for a specified r , K , and $T_{c(C)}$, the receiver should typically be allowed to run at the rate closest to but greater than the ideal bit rate.

If the receiver bit rate is exactly equal to the transmit bit rate of the external transmitting equipment, the sampling of incoming bits will occur at times close to the centers of the corresponding bit intervals. Some error is introduced by the latency between the falling edge of the start bit and the time at which the start bit is detected. The maximum value of that error (e_1) is equal to one period of the timer interrupt.

Additional error is introduced if the receiver bit rate differs from the bit rate of the incoming data stream. Let the bit duration dictated by the timer interrupt rate be denoted by T_1 and let the bit duration of the incoming data be denoted by T_2 . The error introduced by the inequality of T_1 and T_2 (e_2) for the n 'th bit is given by

$$e_{2(n)} = (T_1 - T_2) \times (n - 1/2) \quad (1)$$

The start bit corresponds to $n = 1$. The cumulative error for one word is equal to e_2 evaluated at $n = \text{WORD_LEN}$.

Still another source of error is the latency associated with multicycle instructions. Should a timer interrupt occur during execution of a multicycle instruction or repeat loop, an error e_3 will delay the sampling of $\overline{\text{BIO}}$ by a minimum of zero and a maximum of $I - 1$ cycles, where I is the length (in cycles) of the longest instruction or repeat loop.

The total difference between the sampling time and a corresponding bit interval center is the sum of e_1 , e_2 , and e_3 . In general, the absolute value of the sum of e_1 , e_2 , and e_3 must be less than one-half the duration of one bit in the incoming data stream in order that all sampling instants fall in corresponding bit intervals; i.e.,

$$|e_1 + e_2 + e_3| < T_2/2 \quad (2)$$

The above constraint is appropriate for a receive signal having negligible rise and fall times and equal space and mark durations. If either of these conditions is not satisfied, the constraint expression should be modified accordingly.

Worst-Case Error Analysis

Following are descriptions of the two worst-case scenarios in terms of the three error components. The results of this analysis are then plugged in the constraint expression given in (2) to yield a description of the error constraint in terms of rate difference, K and I .

If the incoming data rate is higher than the receiver bit rate, e_1 , e_2 , and e_3 are all greater than or equal to zero. The worst-case value of e_1 is its maximum value given by

$$\begin{aligned} e_{1(\max)} &= t_{c(C)} \times [\text{PRD reg} > + 1] \\ &= T_1/K \end{aligned}$$

The e_2 contribution is the cumulative error resulting from the inequality of T_1 and T_2 and is given by

$$e_{2(\max)} = (T_1 - T_2) \times (\text{WORD_LEN} - 1/2)$$

The worst-case value of e_3 is given by

$$e_{3(\max)} = (I_{\max} - 1) \times t_{c(C)}$$

If the incoming data rate is lower than the receiver bit rate, e_1 , e_2 , and e_3 are all less than or equal to zero. The worst-case value of e_1 is its minimum value given by

$$e_{1(\min)} = 0$$

The e_2 contribution is the cumulative error resulting from the inequality of T_1 and T_2 and is given by

$$e_{2(\min)} = (T_1 - T_2) \times (\text{WORD_LEN} - 1/2)$$

The worst-case value of e_3 is given by

$$e_{3(\min)} = 0$$

The error constraint (2) is thus satisfied if the following pair of inequalities is satisfied:

$$e_{1(\min)} + e_{2(\min)} + e_{3(\min)} > -T_2/2 \quad (3)$$

$$e_{1(\max)} + e_{2(\max)} + e_{3(\max)} < T_2/2$$

where expressions for the extreme values of each error component are given above.

The inequalities in (3) specify the overall constraint on maximum rate difference, minimum value of K and maximum value of I . For example, suppose

$$\begin{aligned} T_2 &= 0.100 \text{ ms} \\ T_1 &= 0.103 \text{ ms} \\ I &= 20 \\ t_{c(C)} &= 100 \text{ ns} \\ \text{WORD_LEN} &= 10 \end{aligned}$$

Since T_1 is sufficiently close to T_2 , the first inequality in (3) is satisfied:

$$0 + [(103 \times 10^{-6}) - (100 \times 10^{-6})] \times (10 - 0.5) + 0 > (-100 \times 10^{-6})/2$$

Evaluation of the second inequality in (3) yields

$$\begin{aligned} (103 \times 10^{-6})/K + [(103 \times 10^{-6}) - (100 \times 10^{-6})] \times (10 - 0.5) + (20 - 1) \\ \times (100 \times 10^{-9}) < (100 \times 10^{-6})/2 \end{aligned}$$

or

$$K \geq 6$$

Thus (2) translates into a specification for the minimum value of K for a given T_1 , T_2 , I , $t_{c(C)}$, and WORD_LEN .

In summary, considerations must be made with respect to the data rate of the external transmitting equipment, the data rate resulting from the timer interrupt rate, and the latencies associated with start bit detection and multicycle instructions. The two inequalities in (2) must be satisfied for all bits for proper UART operation.

Loopback Test

In the source code given, the XMT and RCV routines are structured to implement a loopback test at 9600 bps, 7 data bits, 1 stop bit and odd parity. The circuit shown in Figure 7 can be used to interface to RS-232-compatible transmit and receive lines. No other RS-232 signals are supported.

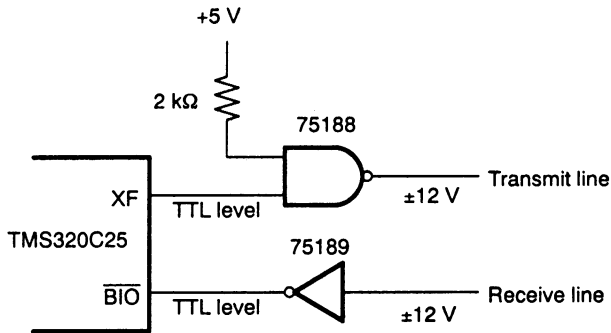


Figure 7. RS-232 Interface

References

- [1] *TMS7000 Family Data Manual*, literature number SPND001C, Texas Instruments, 1989.
- [2] McNamara, John, *Technical Aspects of Data Communications*, Digital Equipment Corporation, 1982.
- [3] *Data Communications Standards*, McGraw-Hill, 1982.
- [4] *Second-Generation TMS320 User's Guide*, literature number SPRU014A, Texas Instruments, 1989.

Appendix

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*****
* SOFTWARE UART USING THE TMS320C25
*
* WRITTEN BY: DAVE ZALAC
*            1/31/89
*
* PROGRAM MEMORY REQUIREMENT: 332 WORDS
* DATA MEMORY REQUIREMENT: 23 WORDS
* MAIN/AIN BIT RATE: 19.2 Kbps
*
*
* ASSEMBLY-TIME CONSTANTS SECTION
*
* STATUS BIT MASKS
*
*****
RFE_MASK .set 001h ; RECEIVE PARITY ERROR
RFE_SET .set 001h
RFE_CLR .set 0FEh
*
FRMR_MASK .set 002h ; RECEIVE FRAMING ERROR
FRMR_SET .set 002h
FRMR_CLR .set 0FDh
*
RPL_MASK .set 004h ; RECEIVE DATA OVERRUN
RPL_SET .set 004h
RPL_CLR .set 0FBh
*
RDL_MASK .set 008h ; RECEIVE DATA AVAILABLE
RDL_SET .set 008h
RDL_CLR .set 0F7h
*
RDE_MASK .set 010h ; RECEIVER ENABLE
RDE_SET .set 010h
RDE_CLR .set 0EFh
*
TDL_MASK .set 020h ; TRANSMIT DATA AVAILABLE
TDL_SET .set 020h
TDL_CLR .set 0DFh
*
*****
* USTAT REGISTER
*
* BIT # 5 4 3 2 1 0
* FLAG TDA RDE RDA RDR RPE
* INITIAL VALUE 0 0 0 0 1 0
*****
*
*****
* PORT ADDRESS OF UART STATUS REGISTER
*
*****
* INITIAL VALUES OF UART PARAMETERS
*
*****
IPACTIVE .set 1 ; IPACTIVE = 1, INDICATES REV PARITY
; CHECKING ON
IKM1 .set 15 ; TINT RATE = K/BIT RATE; KM1 = K-1
IKZM1 .set 7 ; KZM1 = K/2-1
IMK1 .set 6 ; NUMBER OF DATA BITS/MORO
IMORLLEN .set 10 ; TOTAL # OF BITS/MORO
I000 .set 1 ; O00 = 1 SELECTS ODD PARITY
; O00 = 0 SELECTS EVEN PARITY
ITINTPER .set 259 ; TINTPER = TINTER INTERRUPT PERIOD, IN
; CLKOUT1 CYCLES, TINTER SHOULD BE SET TO
; (1/(K*RATE(C)))-1, ROUNDED ROUNDED DOWN
; TO THE NEAREST INTEGER, R IS THE BIT RATE
; IN BPS.
*****
* REGISTER ASSIGNMENTS
*
*****
TAG .set 7 ; DIVIDES TINT RATE BY K (MRT)
RMS .set 6 ; DIVIDES TINT RATE BY K (RCV)
XBITS_REG .set 5 ; USED BY TRANSMITTER TO COUNT DATA BITS
RBITS .set 4 ; USED BY RECEIVER TO COUNT DATA BITS
XMRG .set 3 ; USED BY XCOMPOSE TO COUNT DATA BITS
IMP_PTR .set 2 ; POINTS AT LOCATION TO PUT RECEIVED DATA
OPT_PTR .set 1 ; POINTS AT DATA TO BE TRANSMITTED
*****
* END ASSEMBLY-TIME CONSTANTS SECTION
*
*****
* DATA MEMORY SPACE RESERVATION
*
*****
* VARIABLES
*
*****

```

```

    .bss
    RSTAT,1
    USTAT,1
    SHF,1
    IMP,1
    DMARD,1
    RPARE,1
    RPART,1
    XSTOP,1
    XPAR1,1
    XPAR2,1
    XPAR10G,1
    XDATA,1
    XTSF,1
    XTEMP,1
    RPACTIVE,1
    KRI,1
    KZHI,1
    NHI,1
    WORDLEN,1
    ODD,1
    TIMPER,1
    TIE,1

; RECEIVER STATUS
; UART STATUS REGISTER
; SHIFT FACTOR
; RECEIVED BIT
; RECEIVER SHIFT REGISTER
; LOCALLY-GENERATED PARITY
; INITIAL VALUE OF RPARE
; STOP BIT STRING
; INITIAL VALUE OF XSTOP
; TRANSMIT PARITY
; RECEIVE BUFFER
; TRANSMIT PARITY TOGGLE MASK
; TRANSMIT DATA
; TRANSMIT SHIFT REGISTER
; SCRATCH VARIABLE
; RECEIVE PARITY ACTIVE (0/1)
; K-1
; K/2-1
; N-1
; OVERALL WORD LENGTH
; SELECTS ODD/EVEN PARITY (1/0)
; TIMER INTERRUPT PERIOD
; TRANSMIT BUFFER EMPTY

*****
*
* PROGRAM SECTION
*****
.sect "vectors"
RS B MAIN
.space (RS*24-8)*16; POSITION TINT VECTOR
TINT B XINT ; BRANCH TO XMT/RCV ROUTINE
.text
MAIN LPMK 0 ; INITIALIZE
ZAC ZAC 5 ; NO GLOBAL MEMORY
SDM SDM ; SET OVERFLOW MODE
SPM SPM 0 ; SET SIGN-EXT. MODE
CHFD CHFD ; CONFIGURE BLOCK 80 AS DATA MEMORY
CALL UART_INIT ; INITIALIZE UART
LACK LACK 10N ; ENABLE RECEIVER
USTAT USTAT

*****
; THE USER'S PROGRAM SHOULD APPEAR HERE.
*
* SELF B SELF
*
* END PROGRAM SECTION
*****
; UART INITIALIZATION ROUTINE
*
* THE FOLLOWING CODE INITIALIZES THE UART PER THE VALUES IN THE ASSEMBLY-
* TIME CONSTANTS SECTION ABOVE. ROUTINE PARINIT IS A SUBSET OF UART_INIT
* AND MAY BE CALLED INDEPENDENTLY.
*****
UART_INIT LPMK 0 ; INITIALIZE UART PARAMETER VARIABLES
*
LACK LACK ; IMPACTIVE ; INITIALIZE UART PARAMETER VARIABLES
SACL SACL RPACTIVE
LALK LALK IRTM
SACL SACL KPI
LALK LALK KZHI
SACL SACL KZHI
LALK LALK IRTM
SACL SACL NHI
LALK LALK IMRDLEN
SACL SACL WORDLEN
LACK LACK IODD
SACL SACL ODD
LALK LALK ITIMPER
SACL SACL TIMPER
LACK LACK 1
SACL SACL TIE ; INITIALIZE TIMER PERIOD
ZALS ZALS 2 ; INITIALIZE TIMER PERIOD
SACL SACL 3
LALK LALK OFFCRH ; ENABLE TINT ONLY
SACL SACL 4
ZAC ZAC
SACL SACL RSTAT ; RSTAT = 0
SACL SACL USTAT ; USTAT = 0
OUT OUT USTAT,URTPORT
LAR LAR TMS, KRI ; XMT WAIT: = 1/R
LAR LAR XBITS_REG, WORD_LEN ; INITIALIZE XBITS_REG
ZAC ZAC
SACL SACL DMARD ; DMARD = 0
LALK LALK IMP_PTR, 0200h ; INITIALIZE DATA POINTERS
LALK LALK OPT_PTR, 0200h
*
*****
; CONFIGURE PARITY-RELATED CONSTANTS
*****

```

```

* PHRINIT ZALS
  RPAR SKAL
  XPAR1 SKAL
  LACK 1
  XPAR0S SKAL
  OFFPA LACK
  XSTOP SKAL
  LARP 0
  LAR 0,INT1
  MRR **
  SHIF1 LAC XPAR1,1
  LAC XPAR1
  LAC XPAR0S,1
  SKAL XPAR0S
  LAC XSTOP,1
  SKAL XSTOP
  BMWZ SHIF1

* LAC XSTOP,1
  XSTOP

* EINT
  RET

*****
* END UART_INIT
*****

* TRANSMITTER
*****
* XMIT TBE
  BZ NOT_EMPTY ; IF NOT ZERO THEN EMPTY
*
  ZALS USTAT
  ANBK TDLASK
  BZ OUT ; IF TDR = 0, THEN SKIP TO RECEIVER
*
  ZALS TDR1A
  SKAL TSHF ; TSHF = TDR1A
*
  ZALS USTAT
  ANBK TDLCLR
  SKAL USTAT
  OUT USTAT,UARTPORT
  ZAC
  TBE

* NOT_EMPTY LARP TDR ; IF WAITING, SKIP TO RECEIVER
*****

*
  ZALS RCV
  TSHF ; OUTPUT NEXT BIT
  ROR
  BC XONE
  NOP ; TIME COMPENSATION
  RIF
  B XZERO
  SKF
  SKAL TSHF
  LAR TDR,INT1 ; WAIT 1 BIT INTERVAL BEFORE OUTPUTTING
  ; NEXT BIT
*
  LARP XBITS_REG ; IF LAST BIT, SET THE
  BMWZ RCV
*
  LACK 1 ; SET TBE
  SKAL TBE
*
  LAR XBITS_REG,WORDLEN
*
  CALL INT ; SIGNAL END-OF-WORD
*****
* END TRANSMITTER
*****

* RECEIVER
*****
* ZALS USTAT ; SKIP RECEIVE ROUTINE IF RDE = 0
  ANBK RDE_LSK
  BMWZ CONT
  B RETURN
*
  CONT ZALS RSTAT ; RSTAT = 0 IMPLIES WAITING FOR START BIT
  BMWZ NOTSTART
*****
*
  RSTAT = 0
*****
*
  BIOZ STARTBIT ; LOOK FOR START BIT ON BIO/
  B RETURN
*****
*
  STARTBIT LACK 1 ; UPDATE RSTAT
  SKAL RSTAT
  LAR RMR,CONT1 ; WAIT 1/2 BAUD INTERVAL AFTER START BIT
*****

```



```

* RCVPAR_OK_LAR RMG_INTI ; WAIT 1 BRND INTERVAL
*
* LACK 4
* SACL
*
* B RSTAT
* RETURN
*
*****
* RSTAT = 4
*
*****
* NOTPAR BLDZ STOPO ; INVALID STOP BIT
*
* B STOP_OK
*
* STOPO ZALS USTAT ; REPORT FRAMING ERROR
* ORK FMRBR_SET
* SACL USTAT
* OUT USTAT,UARTPORT
*
* STOP_OK ZALS USTAT ; OVERRUN DETECT
* ANBK ROLANSK
* BL MLOVERRUN
*
* OVERRUN ZALS USTAT ; REPORT OVERRUN
* ORK ROL_SET
* SACL USTAT
* OUT USTAT,UARTPORT
*
* MLOVERRUN ZALS DMARD ; COPY DATA TO RDATA
* SACL RDATA
*
* ZALS USTAT ; SET RDA
* ORK ROL_SET
* SACL USTAT
* OUT USTAT,UARTPORT
*
* ZAC
* SACL RSTAT
* SACL DMARD
*
* ZALS RPARI
* SACL RPAR
*
* CALL RECV
*
* RETURN EINT
* RET
*
*****
*
*****
* END RECEIVER
*
*****
* TRANSMITTER DATA FEED ROUTINE
*
* THE FOLLOWING INT/RCV ROUTINES IMPLEMENT A LOOPBACK TEST.
*
*****
* INT RET
*
*****
* END INT
*
*****
* RECEIVER DATA RECOVERY ROUTINE
*
* THIS SET OF INT/RCV ROUTINES IMPLEMENTS A LOOPBACK TEST.
*
*****
* RECV CALL GET_DATA ; GET RECEIVED DATA
* CALL XCOMPOSE ; REFORMAT
* CALL PUT_DATA ; TRANSMIT SAME DATA
* RET
*
*****
* END RCV
*
*****
* TRANSMITTER INTERFACE ROUTINE
*
* PUT_DATA COPIES CAR(OPT_P1R1) TO TDATA AND SET TDA
*
*****
* PUT_DATA LAMP OPT_P1R1 ; COPY DATA TO TDATA
* SACL *
* SACL TDATA
*
* ZALS USTAT ; SET AND REPORT TDA
* ORK TDA.SET
* SACL USTAT
* OUT USTAT,UARTPORT
*
* RET
*
*****

```



```

* END PUT_DATA
*****
* RECEIVER INTERFACE ROUTINE
*
* GET_DATA COPIES DATA TO LOCATION POINTED AT BY @R(INP_PTR) AND CLEARS
* RDA, RDR, FRH, AND RPE
*****
*
* GET_DATA LARP INP_PTR ; COPY DATA
* ZALS RDATA
* SACL *
*
* ZALS USTAT
* ANDK RDA_CLR ; CLEAR RDA
* ANDK OFFERR ; CLEAR ERROR FLAGS
* SACL USTAT
* OUT USTAT,UARTPORT
*
* RET
*****
*
* END GET_DATA
*****
* TRANSMIT DATA COMPOSE ROUTINE
*
* XCOMPOSE ADDS START, STOP, AND PARITY BITS TO DATA POINTED AT BY
* @R(OPT_PTR)
*****
*
* XCOMPOSE ZALS XPAR1 ; COMPOSE TRANSMIT WORD
* SACL XPAR ; INPUT IS @R(OPT_PTR)
* ; OUTPUT IS @R(OPT_PTR)
*
* LARP OPT_PTR
* ZALS *
* SACL TEMP
*
* LAR XBREG, NMI1 ; NUMBER OF TRANSMIT DATA BITS
* LARP XBREG
*
* ZALS TEMP ; DETERMINE PARITY BIT VALUE
* ROR ROR
* SACL TEMP ; NO PARITY CHANGE IF SPACE
* BNC INDCORR
*
* C_PAR ZALS XPAR ; TOGGLE PARITY-GEN BIT IF MARK
* XOR XPAR10G

```

```

* SACL XPAR
*
* INDCORR BANC C_PAR
*
* LARP OPT_PTR
* LAC *,1 ; ADD START BIT
* XPAR ; ADD PARITY BIT
* OR XSTOP ; ADD STOP BITS
* SACL *
*
* RET
*****
*
* END XCOMPOSE
*****

```